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ON USING EQUALITY-CONSTRAINT ALGORITHMS
FOR INEQUALITY CONSTRAINED PROBLEMS[†]

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Consider the inequality constrained problem

$$P_0: \quad \max f(x), \text{ subject to} \\ g_i(x) \leq 0, \quad i = 1, \dots, m$$

where all functions are real valued and continuous on R^n . Suppose x^* is an optimal solution to P_0 and define $I = \{i : g_i(x^*) = 0\}$. Recently a number of references have appeared (see, for example, [2,4]) which discuss the possibility that certain techniques for solving equality constrained problems can be effectively extended to problems like P_0 by first identifying I and then solving the problem

$$\hat{P}: \quad \max f(x), \text{ subject to} \\ g_i(x) = 0, \quad i \in I.$$

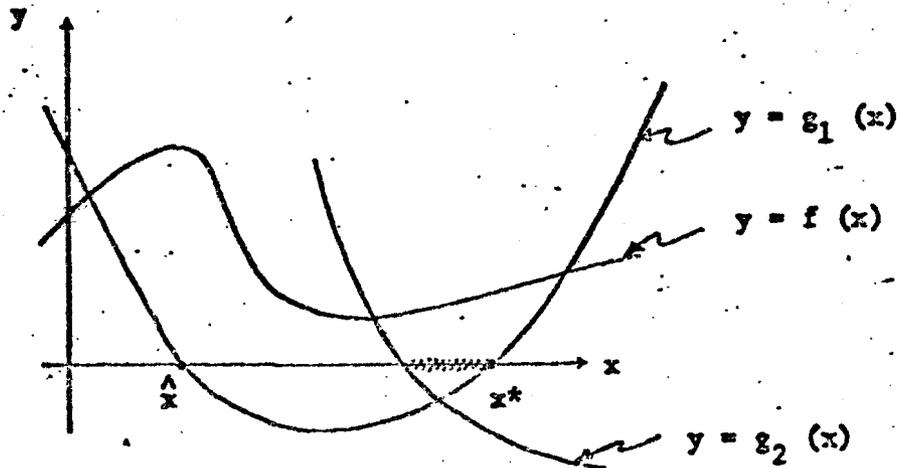
There has also been some evidence of computational success in implementing such techniques and solving an assortment of "standard" test problems [3]. The purpose of this note is to establish conditions on problem P_0 which make its replacement by \hat{P} valid, but also to present theoretic difficulties which must inevitably be reckoned with by those researchers who would wish to elaborate the development of these techniques. The difficulties

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we refer to are due mainly to the facts that (i) a concave problem P_0 in general does not produce a convex constraint set in \hat{P} , and (ii) known penalty functions for \hat{P} , in the case when f is concave and the constraint functions convex, are not themselves concave functions. This point will be further discussed below. Prior to that, we first wish to note that there appears to be a not uncommon impression that P_0 , at least in theory, can always be replaced with \hat{P} . The following example shows clearly that this is generally not true.



Here the constraint set $S_0 = \{x : g_i(x) \leq 0, i = 1, 2\}$ is shaded, x^* solves P_0 , $I = \{1\}$, and \hat{x} solves \hat{P} . The following lemma gives a set of conditions under which, in theory, the replacement of P_0 with P is valid.

Lemma: Suppose x^* is a global solution to P_0 , f is strictly quasi-concave¹

¹ f is strictly quasi-concave on R^n if for each pair of points $x^1, x^2 \in R^n$, $f(x^1) > f(x^2)$ implies $f(tx^1 + (1-t)x^2) > f(x^2)$, $0 < t < 1$.

on R^n , and the set $S_1 = \{x: g_1(x) \leq 0\}$ is convex for each $i \in I$.¹ Then x^* is a global solution to \hat{P} .

Proof: Define $S_0^I = \{x: g_i(x) \leq 0, i \in I\}$; suppose $\hat{x} \in S_0^I$ but $f(\hat{x}) > f(x^*)$. Define $x_t = tx + (1-t)x^*$, $0 \leq t \leq 1$. Suppose $t > 0$; by strict quasiconcavity of f , $f(x_t) > f(x^*)$. Thus since x^* solves P_0 , it follows that $x_t \notin S_0$, $0 < t \leq 1$.

Since S_1 is convex for each $i \in I$, and $\hat{x}, x^* \in S_0^I$, the line segment connecting x^* and \hat{x} is contained in S_0^I . But we will show that there is a $t \in (0, 1]$ such that $x_t \notin S_0^I$, yielding a contradiction. For $i \notin I$, $g_i(x^*) < 0$. By continuity of g_i 's, there is a $t_1 \in (0, 1]$ such that $0 < t \leq t_1$ implies $g_i(x_t) < 0$, $i \notin I$. Now since $t \in (0, 1]$ implies $x_t \notin S_0$, it must be that for some fixed $t \in (0, t_1]$, there is an index $i_0 \in I$ such that $g_{i_0}(x_t) > 0$. Thus, $x_t \notin S_0^I$. This contradiction shows that $\hat{x} \in S_0^I$ implies $f(\hat{x}) \leq f(x^*)$. Thus, x^* is a global solution to $\max f(x)$, subject to $x \in S_0^I$, and hence a global solution to \hat{P} . \square

Thus, in particular, the replacement notion is theoretically valid for a concave program, P_0 . However, it is important to note that concavity of problem P_0 is not, in general, preserved in problem \hat{P} . That is, the constraint set for \hat{P} need not be concave. Part of the impact of this fact is reflected in the second of the following easily verified assertions for general (not necessarily concave) problems P_0 :

1) If \hat{x} is a global solution to \hat{P} and \hat{x} is feasible in P_0 , then

\hat{x} is a global solution to P_0 .^{2,3}

1 Note that these assumptions imply neither uniqueness of x^* nor convexity of S_0 . Indeed P_0 could have local solutions that are not global.

2 Note that we here "assume away" the important problem of identifying I . Typical procedures for identifying I , for problems with nonlinear constraints, are ad hoc and heuristic, involving the inclusion of "blocking" constraints, the deletion of constraints with a negative Lagrange multiplier, or the utilization of penalty function maximizations to obtain points "close" to x^* so that an educated estimate of I might be made. However, these procedures in general entail scaling difficulties, and they may identify a basis which is not consistent for \hat{P} . It would appear that the problem of identifying I is itself one of substantial difficulty.

3 Bazaraa and Goode have shown in [1] that if \hat{x} is a global solution to P and if P_0 is a strictly concave program, then \hat{x} will in fact be feasible in P_0 .

ii) If x^* is a local solution to P_0 , then x^* is a local solution to the corresponding¹ problem \hat{P} but not conversely.² Specifically, suppose J is some specified subset of $\{1, \dots, m\}$, and suppose \hat{x} is a local solution to the problem $\max f(x)$, subject to $g_i(x) = 0, i \in J$; then, even in the case when P_0 is a concave program, \hat{x} need not be feasible in P_0 ; and even if \hat{x} is feasible in P_0 , it does not follow that it is a local solution to P_0 .

This holds even when $J = I$.

According to ii), even when P_0 is a concave program, one may obtain local solutions to \hat{P} which are not local solutions to P_0 . Furthermore in the case of a concave program P_0 , when replacement by \hat{P} is valid, typical penalty functions for equality constrained problems, as presented in [1] and [3], are not concave. In general it would seem obviously desirable to employ solution techniques which preserve and exploit as much of the useful structure of the original problem as possible. If f is concave and the g_i 's are convex then we can replace P_0 by either of the following problems.

\bar{P} :

$\max f(x)$, subject to

$$g_i(x) \leq 0, i \in I$$

\hat{P} :

$\max f(x)$, subject to

$$g_i(x) = 0, i \in I$$

The replacement is valid in the sense that x^* , a global solution to P_0 , is also a global solution to \bar{P} and \hat{P} . If we employ a typical penalty function for equality constrained problems in solving \hat{P} , e.g.

$$P(x, \delta) = f(x) - \delta \sum_{i \in I} (g_i(x))^2$$

¹We employ the term, corresponding, in order to emphasize the fact that problem \hat{P} depends on x^* .

²Since this assertion applies to each local solution to P_0 , it might, as suggested in a private communication by Professor Luenberger, be used to derive necessary differential optimality conditions for P_0 . That is, such conditions are based only on local properties.

$P(x, \delta)$ is not concave in x . Thus we have lost the nice structure of problem P_0 . Given a parameter sequence $\{\delta_n\}$, $\delta \rightarrow +\infty$, a sequence of global maximizers of P will converge to x^* under fairly weak additional restrictions. However, the non-concavity of \hat{P} raises the difficult practical problem of local maxima of P . All penalty functions we are aware of that are specific for equality constrained problems (i.e. they work for equality constraints but not inequality constraints) have this property that they do not preserve concavity. By comparison, we can choose penalty functions for \bar{P} which retain the very valuable concave structure of that problem. Thus, particularly in the concave case, it seems reasonable that if the original constraints were formulated as inequalities, one might wish to use a penalty function appropriate for inequality constraints (that is, if penalty functions are to be used at all).

Finally, along a slightly different vein, it has been observed by Professor Mangasarian in a private communication that the following problem is always equivalent to P_0 (regardless of convexity).

$$\begin{aligned} \hat{P} \quad & \max f(x), \text{ subject to} \\ & g_i(x) = 0, \quad i \in I \\ & g_i(x) < 0, \quad i \notin I \end{aligned}$$

and that in reality it is this problem which one tends to solve with equality constraint procedures, taking special precaution to satisfy $g_i(x) < 0, i \notin I$. Typically this would be done as follows. Suppose x_k is a current iterate (trial solution) for \hat{P} and I_k is the current guess at I . In determining the next estimate, x_{k+1} , suppose, for some i such that $g_i(x_k) < 0$, the algorithm is tending to force g_i positive e.g. a line search is crossing the contour $g_i(x) = 0$. Then feasibility would be maintained by adding this

"blocking constraint" to the basis I_k . A rule such as this, introduced to maintain feasibility, might retard convergence, particularly if it leads to a zig-zag phenomenon in which one constraint is repeatedly added and deleted from the basis. More explicitly, considering this same point, suppose P_0 is a concave program and one sets out to solve \hat{P} with an equality constraint technique. Even if feasibility is maintained, as just discussed, assertion (ii) above indicates the possibility of converging to a point (a local solution to \hat{P}) which is not a solution to P_0 . Provided P_0 is a concave program, such a false optimum will be indicated by a negative Lagrange multiplier. As first suggested by Rosen [5] in his Gradient Projection method, the algorithm can be continued by deleting any constraint with a negative multiplier, thereby effecting a change of basis. However, the nonconcavity of \hat{P} , resulting in local solutions which are not solutions to P_0 , may in general lead to a waste of time in converging to wrong local optima of \hat{P} . Finally, if P_0 is not concave, then there may be Kuhn Tucker points which are not solutions, but the problem \hat{P} should be expected to have more such points than P_0 . In such cases, all one can hope for is convergence to a Kuhn Tucker point, and hence, in the nonconcave case, it should be more hazardous to try to solve \hat{P} .

In summary, then, even if P_0 is concave, \hat{P} is not concave. Hence, one may have difficulties in determining global solutions to \hat{P} . If the initial starting point is close enough to x^* , and if feasibility is maintained, with anti-zig-zagging techniques, perhaps the difficulty in convergence to x^* will be minimal. However, if the starting point is more remote from x^* , theory would expect a convergence to local solutions

of \hat{P} which are not a solution to x^* . If P_0 is concave, false optima can be detected, so the issue is one of efficiency. If P_0 is not concave, the issue is one of reliability.

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